

# **“The Integration Platform”**

- ◆ **A weak position given the new Windows contract**
  - **IBM’s access to Windows technology ends with what’s in process in September 1993**
- ◆ **IBM only has OEM rights to Windows NT**
  - **They can sell the product or port it to the RS6000**
  - **They CANNOT incorporate Windows NT technology into OS/2**
- ◆ **IBM has no rights whatsoever to the Win32 API or Win32s**

# **Advantages of Windows NT Capability**

- ◆ **Security**
- ◆ **Portability to RISC**
- ◆ **Multiprocessor support**

**IBM can't fix these problems  
with the existing OS/2 code  
base**

- **They must move their  
customers to a new  
operating system (Mach?)**

# **Advantages of Windows NT**

## **Seamless Compatibility**

- ◆ **Windows NT seamlessly runs MS-DOS and Windows apps - including Win 3.1**
  - **One set of printer drivers and fonts**
  - **One place to configure the system**
  - **Single command shell runs MS-DOS, Win16, Win32, OS/2, and Posix apps**

# **OS/2 Compatibility**

## **This is Seamless Windows?**

- ◆ **OS/2 is 2 operating systems pasted together**
  - **2 sets of printer drivers and fonts**
  - **2 places to configure the system**
  - **Special configuration needed to run Windows apps**
  - **Separate DOS and OS/2 command prompts**

# Advantages of Windows NT

## Migrating Apps to 32-bit

- ◆ Win32 is extremely compatible with Win16
  - Apps are easy to port to 32-bit
  - Single 16/32-bit source code can be maintained
  - It's happening!
- ◆ Presentation Manager is different in every detail with Win16
  - Incredibly difficult to move Windows apps to Presentation Manager
  - Very few PM apps (16 or 32-bit)

# Windows NT Architecture

**“True 32”**

**Windows**

**NT**

32-bit APIs		16-bit APIs	
Graphics 32-bit		Window Mgr 32-bit	
NT Kernel 32-bit			
NTFS 32-bit	FAT 32-bit	HPFS 32-bit	
Device Drivers 32-bit			

# OS/2 2.0 Architecture

## “Blue 32”

OS/2

32-bit APIs	16-bit APIs
Graphics 16-bit	Window Mgr 16-bit
OS/2 32-bit	
FAT 32-bit	HPFS 16-bit
Device Drivers 16-bit	

# **Implications of 16-bit code in 2.0**

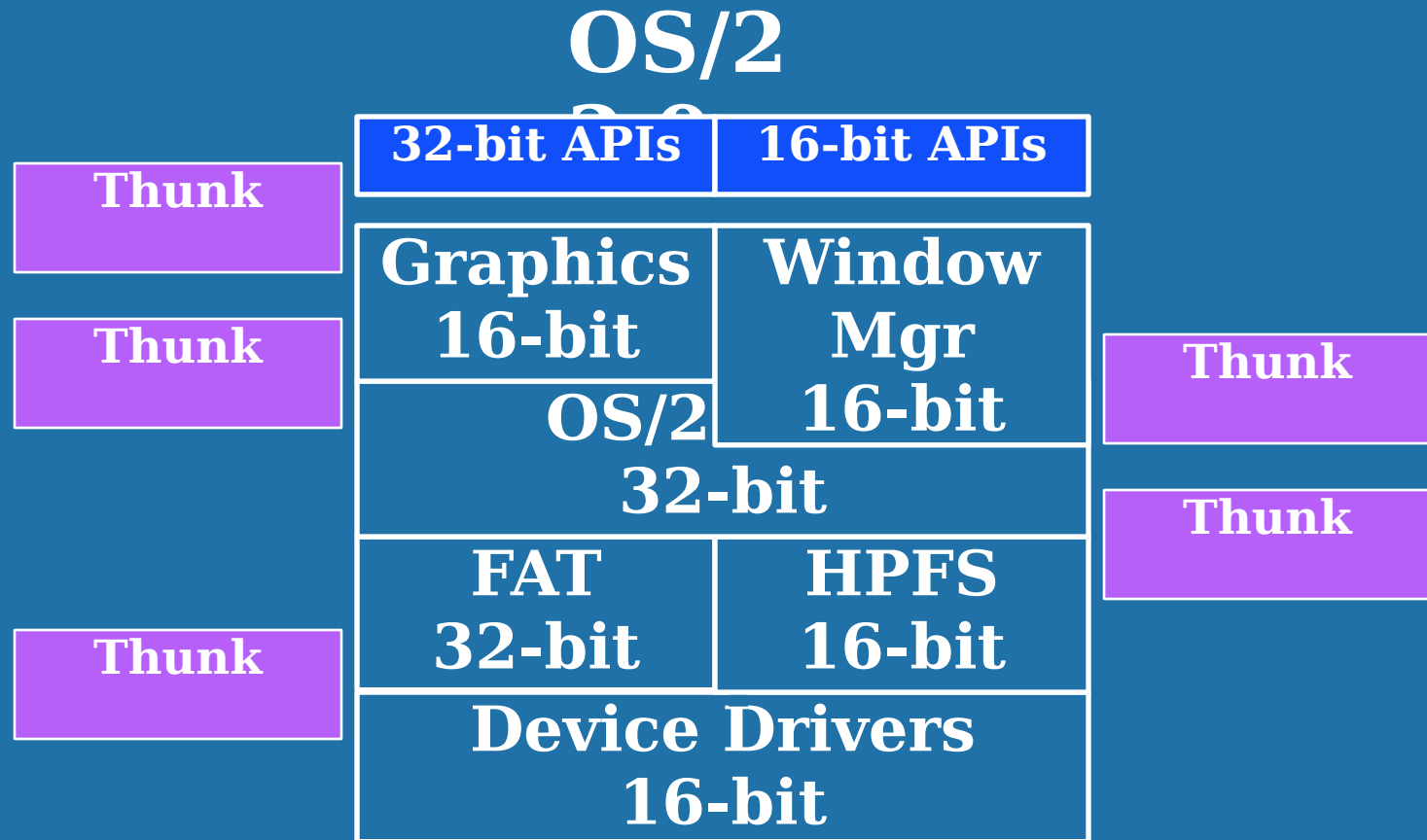
## **Thunking**

- ◆ **Thunks convert 16-bit addresses to 32-bit**
  - **Or 32-bit to 16-bit**
- ◆ **Thunks are pure overhead**
- ◆ **Thunks are required in any OS which supports both 16-bit and 32-bit applications**

**The less thunks the better!**



# Thinking in OS/2 2.0



# Thinking in Windows NT

## Windows

NT

32-bit APIs		16-bit APIs	
Graphics 32-bit		Window Mgr 32-bit	
NT K			
32-bit			
NTFS 32-bit	FAT 32-bit	HPFS 32-bit	
Device Drivers 32-bit			

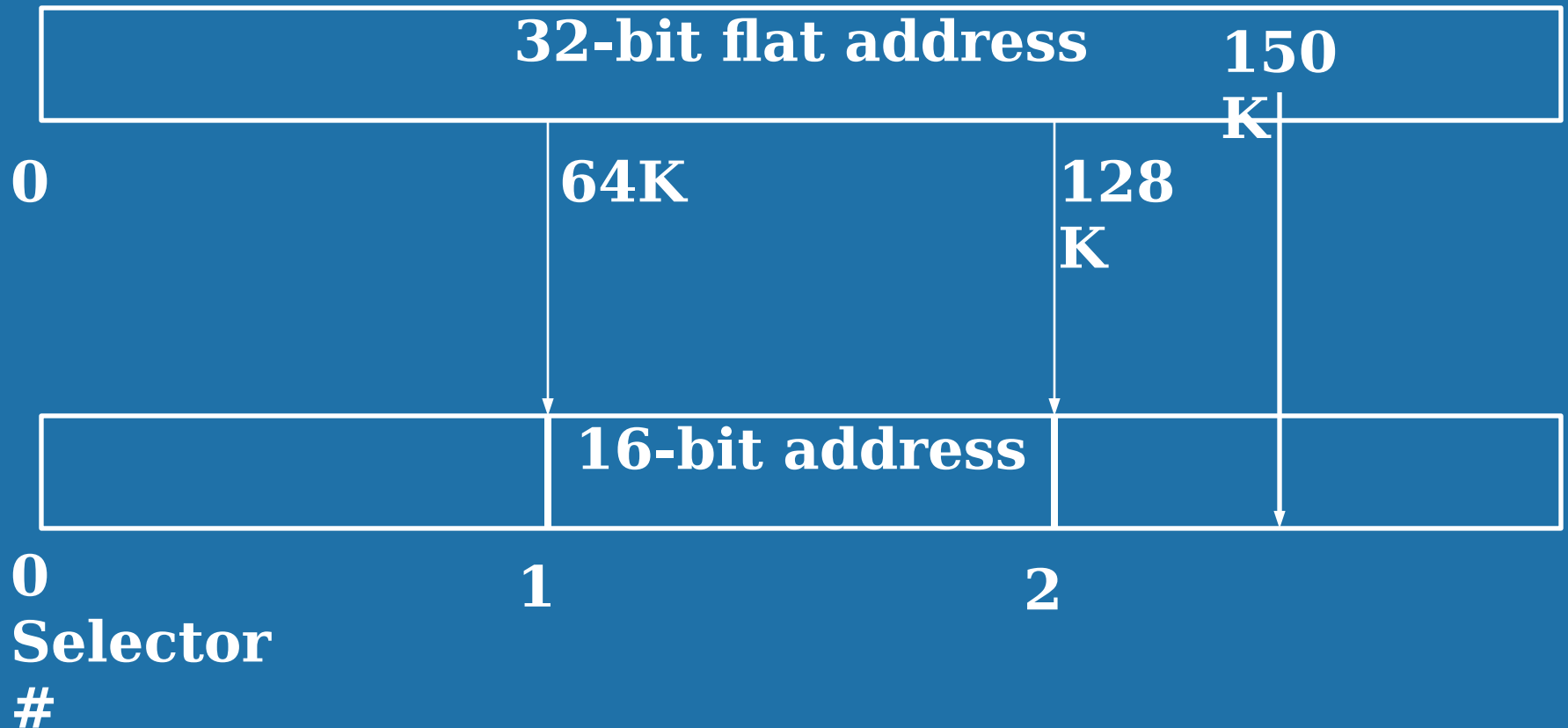
Thunk

# Implications of 16-bit code in 2.0

## Tiling

- ◆ Fast 16/32-bit address conversion in OS/2 is critical because of all the 16-bit code
- ◆ Tiling provides a convenient way to convert 32-bit addresses to 16-bit
  - A 32-bit address is flat - 0:32
  - A 16-bit address contains selectors and offsets - 16:16
- ◆ With tiling, each selector begins on a 64K boundary

# Tiling in OS/2 2.0



**A 32-bit address of 150K = selector 2, offset 22K**

# Downside of Tiling

- ◆ Every data segment begins on a 64K boundary
  - This implies a new 4K page for each data segment
- ◆ Most 16-bit applications create lots of small data segments
  - Efficient way to manage memory on a 16-bit OS
- ◆ Result: Program bloat thru wasted memory
  - If an app allocates 100 1K segments, OS/2 2.0 wastes 300K

**Windows NT does not tile memory**

# **Implications of 16-bit code in 2.0 Arbitrary Limits**

- ◆ **Two types of addressing in 16-bit code:**
  - **Near: 1 data segment (FAST)**
  - **Far: multiple data segments (SLOWER)**
- ◆ **Near addressing limits data size to 64K**
  - **This creates a seemingly arbitrary limit**

# **No Limits with Windows NT**

- ◆ **32-bit addresses are always near**
  - **Segment size = 4GB**
- ◆ **32-bit design of Windows NT eliminates limits with no performance penalty**
- ◆ **Windows NT eliminates limits in other ways:**
  - **64-bit volume and file addressing**
    - **max file size = 20 quintillion bytes**

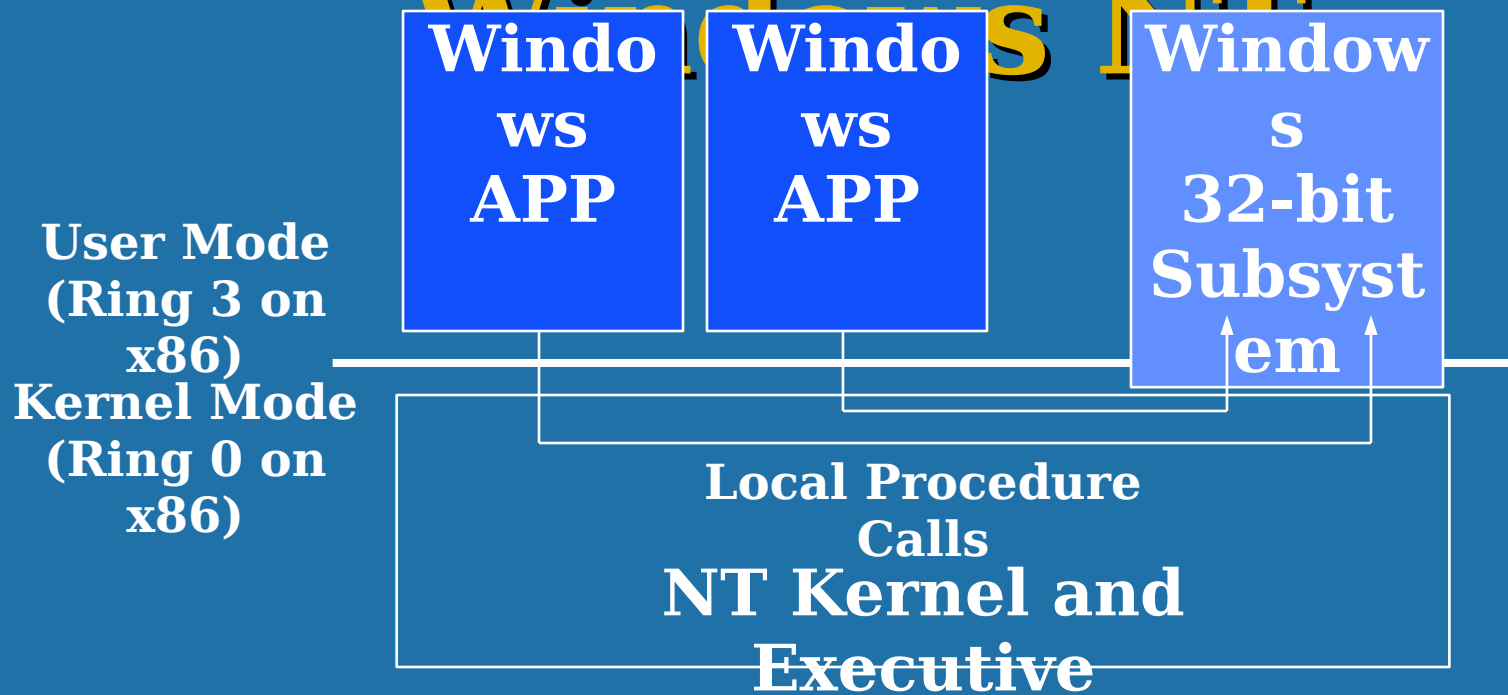
# Address Space Separation in OS/2 2.0



- ◆ Character and server apps run in separate address spaces
- ◆ PM apps share address space with PM
  - A PM app can corrupt PM and crash



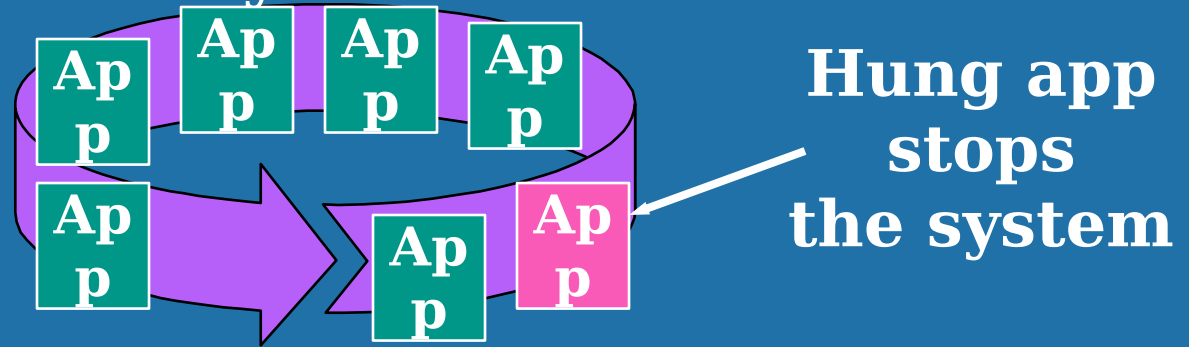
# Address Space Separation in Windows NT



- ◆ 32-bit apps run in separate address spaces
  - Can't crash the system
- ◆ Win16 apps run in 1 address space, but they still can't crash

# Single Message Queue in OS/2

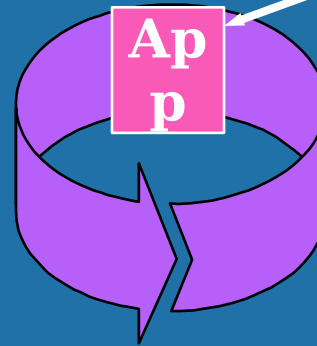
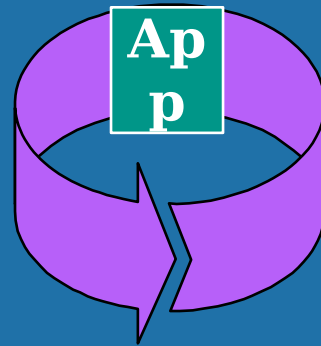
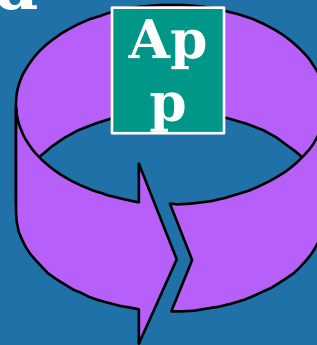
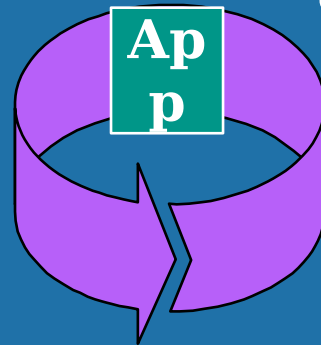
One Queue for the Whole  
System



- ◆ All OS/2 apps share 1 message queue
- ◆ If an app hangs, the system is frozen until a timeout expires
  - The user can't do other work

# Desynchronized Message Queue in Windows NT

One Queue for each  
thread

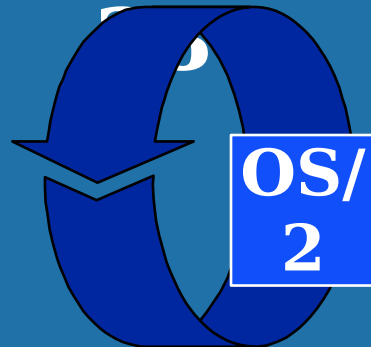


Hung app  
has no  
effect  
on the  
system

- ◆ With Windows NT, it is always possible to switch away from a busy or hung app

# Network Server on OS/2 2.0

- ◆ IBM has not shipped a high performance Lan Server for OS/2 2.0
  - This requires HPFS 386, which is not yet supported on 2.0
- ◆ HPFS 386 bypasses OS/2 scheduling
  - HPFS 386 actually schedules OS/2 as a task



# **Network Server on Windows NT**

- ◆ **Built in!**
  - **No connection limits**
- ◆ **Uses Windows NT Kernel scheduling**
  - **Basis for multi-processor support**

# Other Features in Windows NT

- ◆ One button printer setup
- ◆ Single user account/password in an enterprise (with Lan Manager)
- ◆ FT: Mirroring, striping, striping with parity (with Lan Manager)
- ◆ Built-in backup utility
- ◆ Performance Monitor
- ◆ Unicode
- ◆ Configuration Registry...

# Summary

- ◆ IBM will fight to the death with OS/2
  - Very, very scary
- ◆ Windows keeps moving forward
  - ISV support has shifted to Win32
- ◆ Windows NT is solid and is available **now** to developers!
- ◆ Windows NT blows away OS/2 for corporate computing